

# GRAPH TRAVERSAL ALGORITHM: DEPTH FIRST SEARCH (DFS)

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# WHY STUDY GRAPH TRAVERSAL ALGORITHMS?

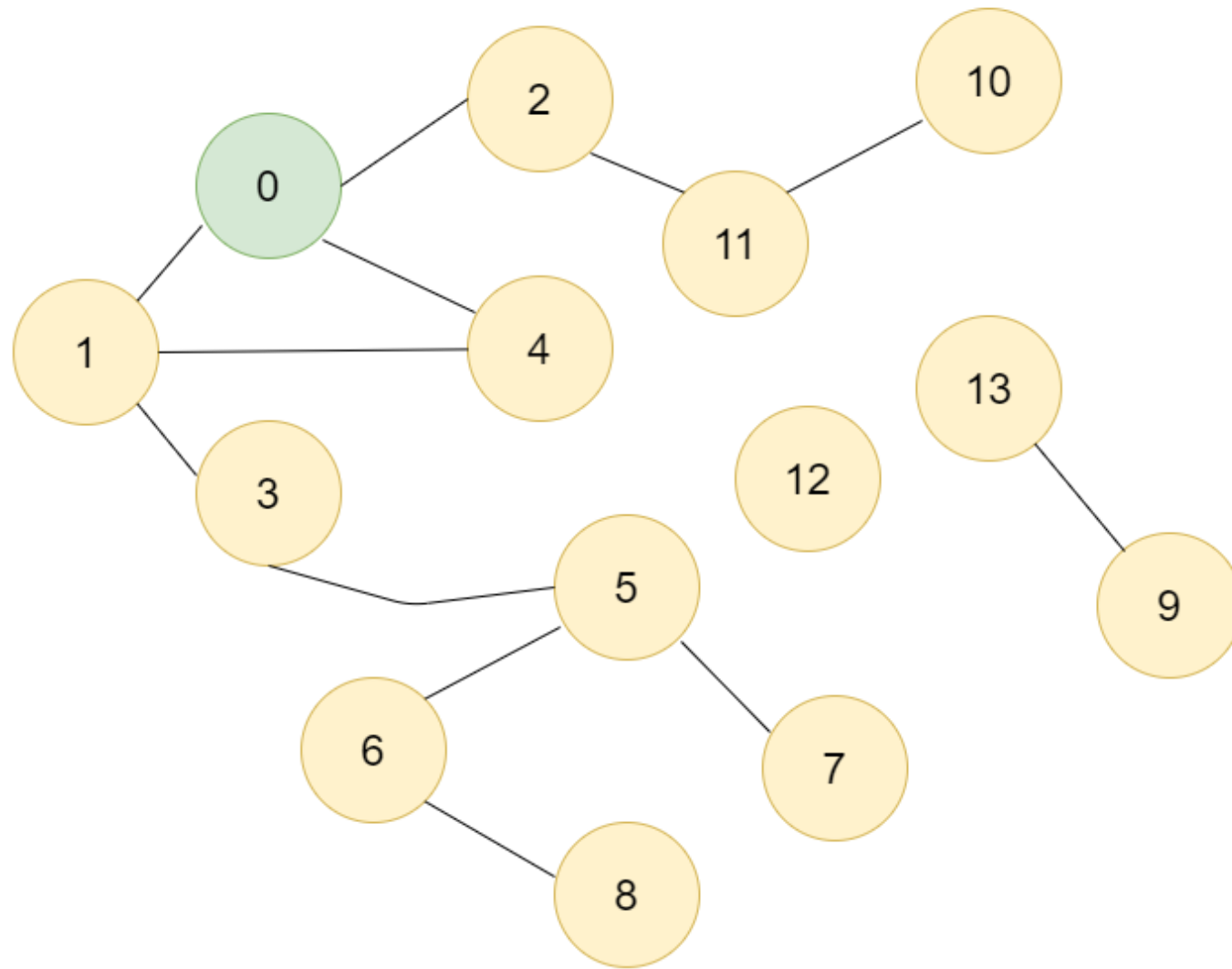
- Graph Reachability Problems
- Cycle Detection
- Topological Sort
- Connected Components
- Flood Fill Problems
- Shortest path in unweighted Graphs
- Strongly connected components (Tarjan's)
  - Only for directed graphs.
- Biconnected Components
- Finding path in an unweighted graph
- Bipartite graph Check
- Articulation Points / Bridges

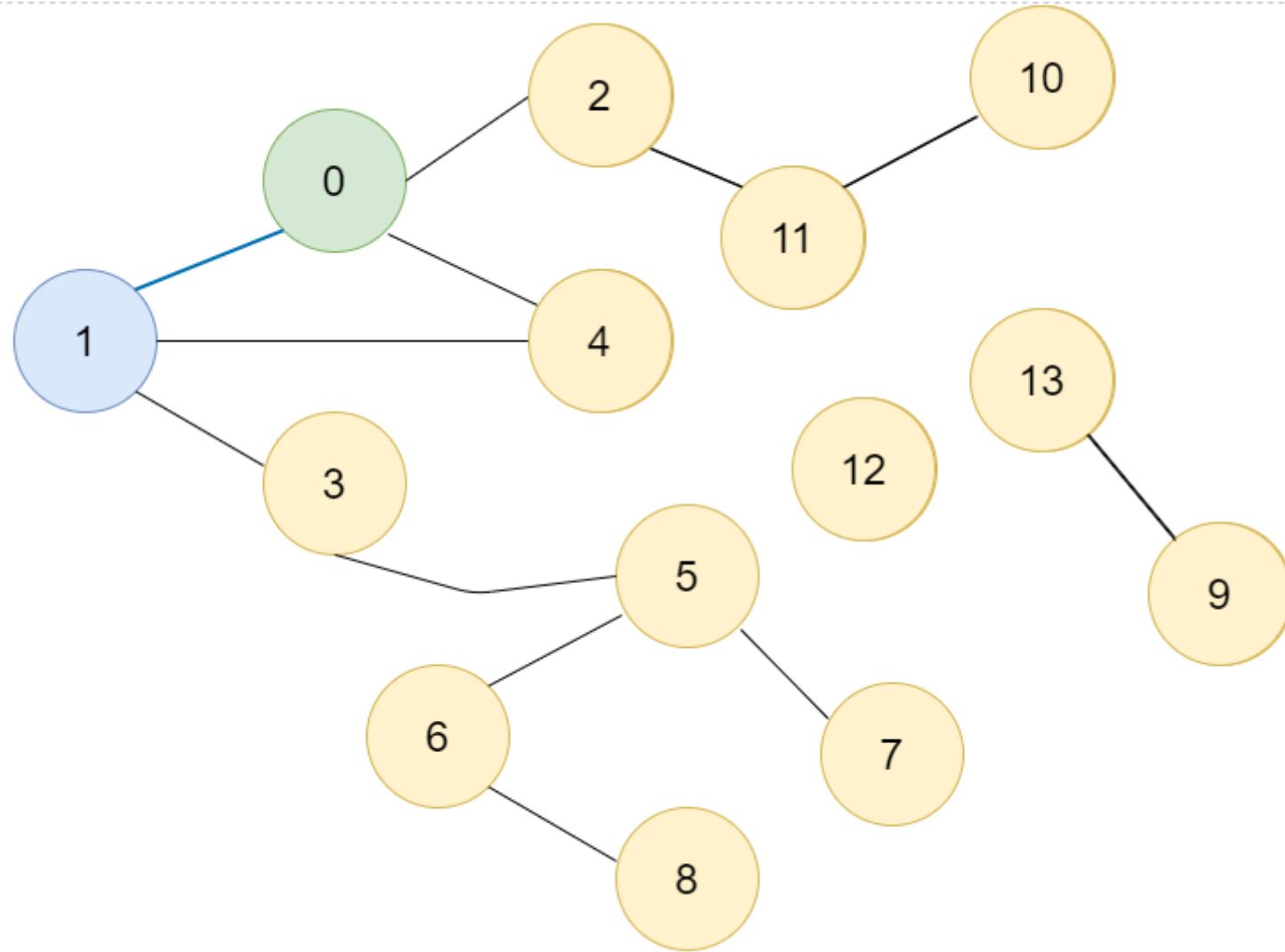
# GRAPH TRAVERSAL ALGORITHMS

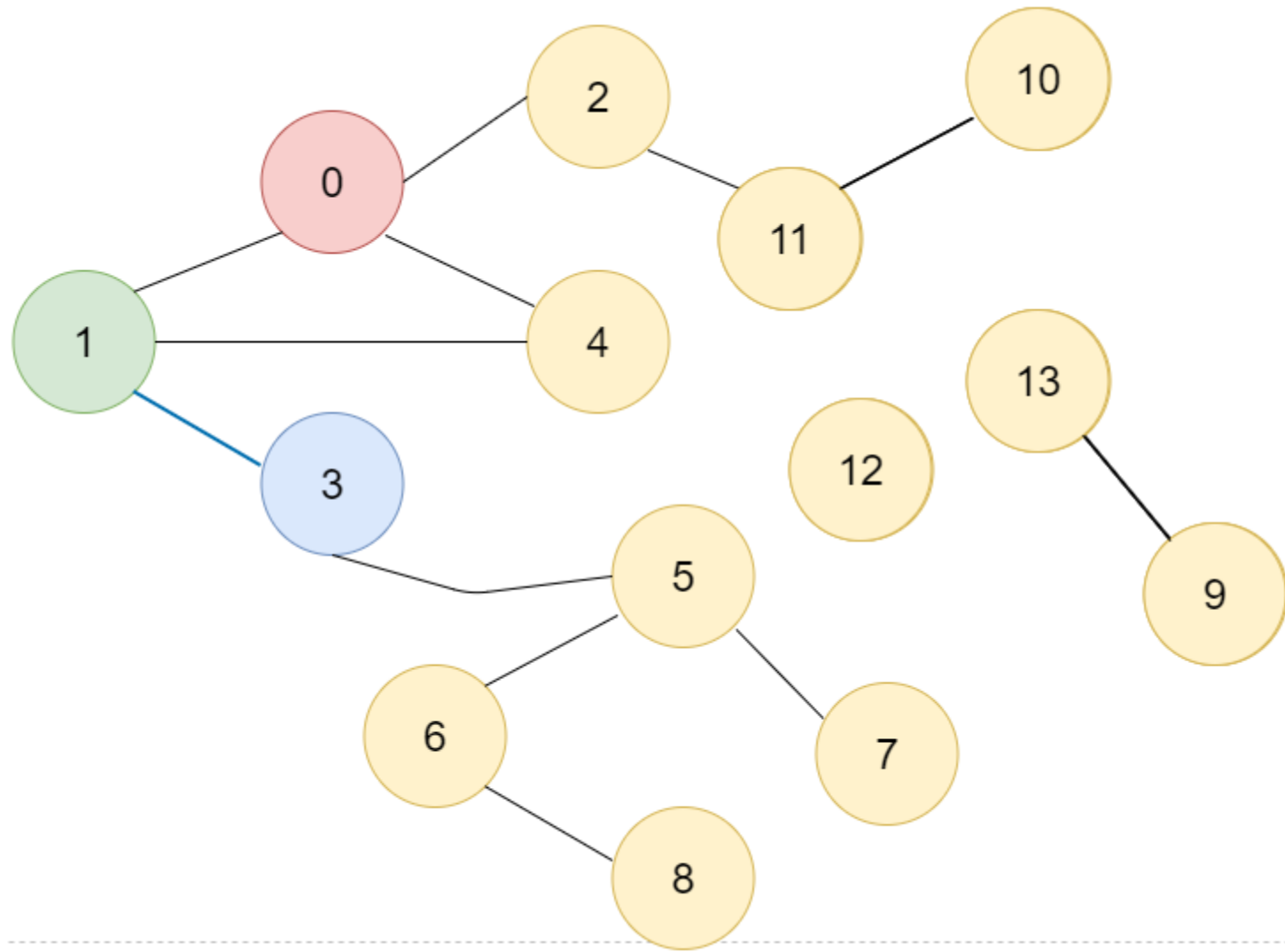
- Depth First Search (DFS)
- Breadth First Search (BFS)

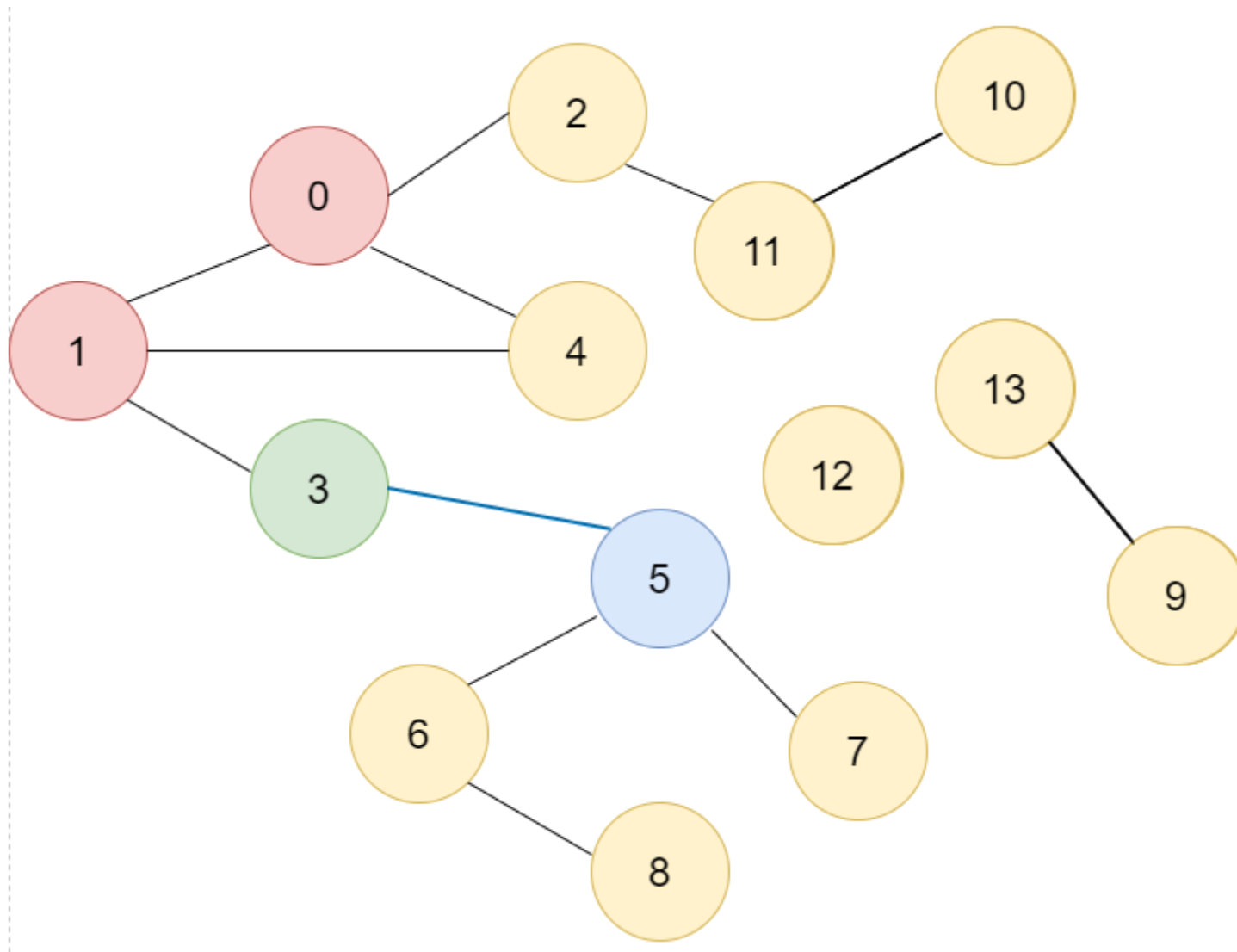
# HOW DEPTH FIRST SEARCH (DFS) WORK?

- ❑ Start from a single node and traverse the graph in a depth first fashion.
  - ❑ Do depth first search at the adjacent nodes.
  - ❑ Keep going deeper until reaches a vertex which does not have any unvisited vertex.
  - ❑ Then backtrack (return) to the previous node and explore other branches/paths from that node.
- ❑ Follows a single path as long as it finds a new node.
  - ❑ Explores a branch completely before moving on to a new branch of the graph
  - ❑ It goes deep first and branches later.
- ❑ Keep tracks of the visited nodes so that no nodes get processed more than once.

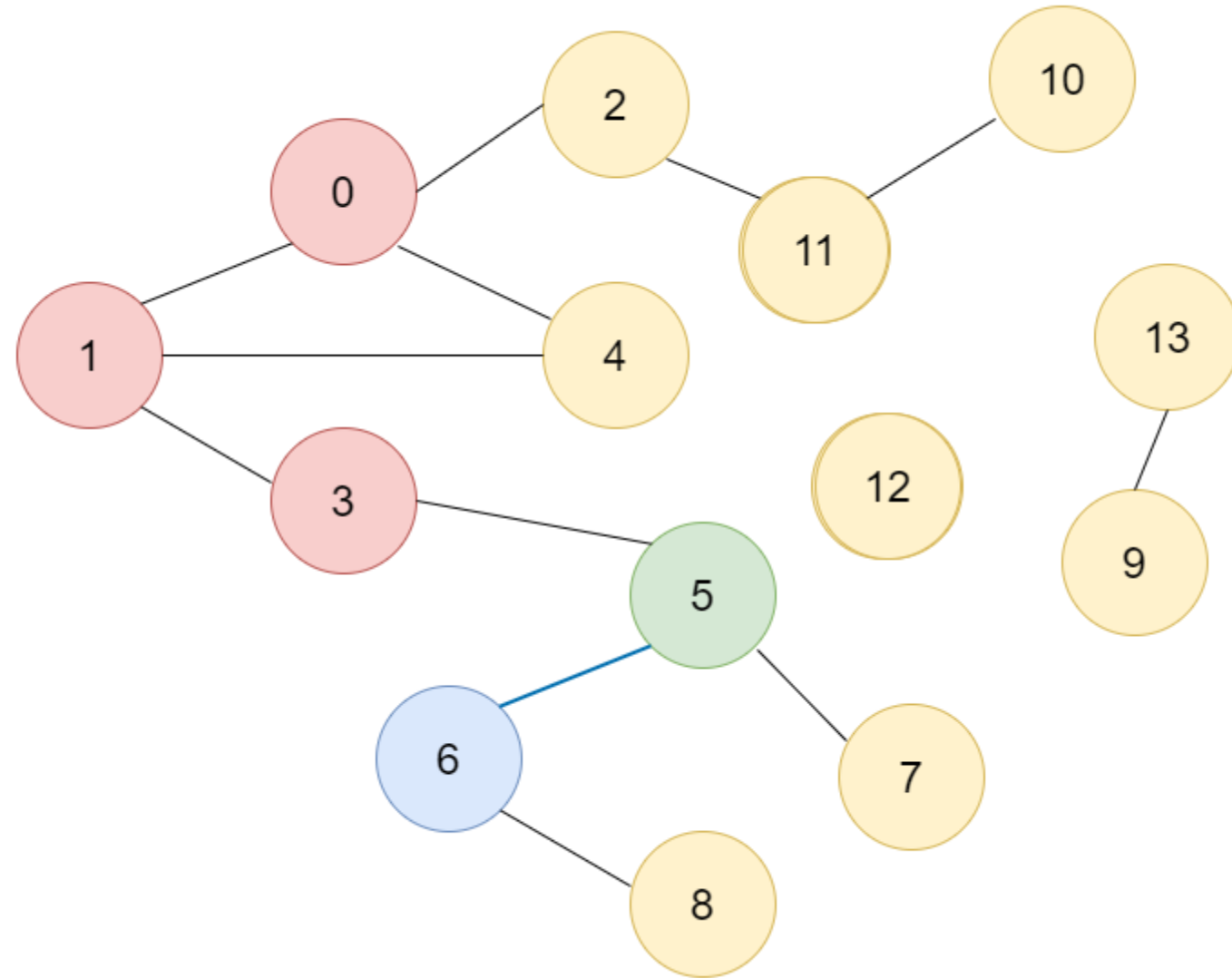


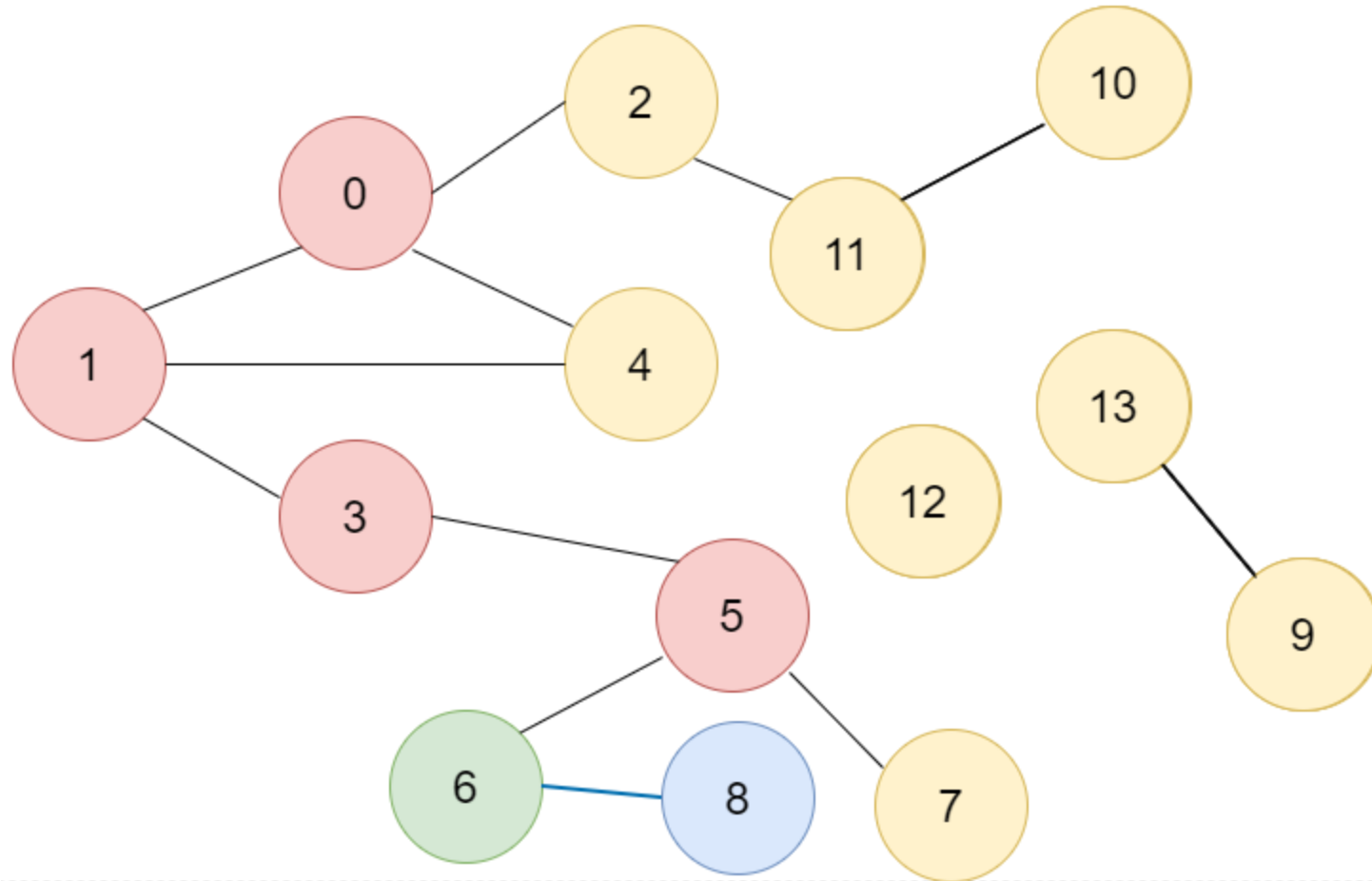


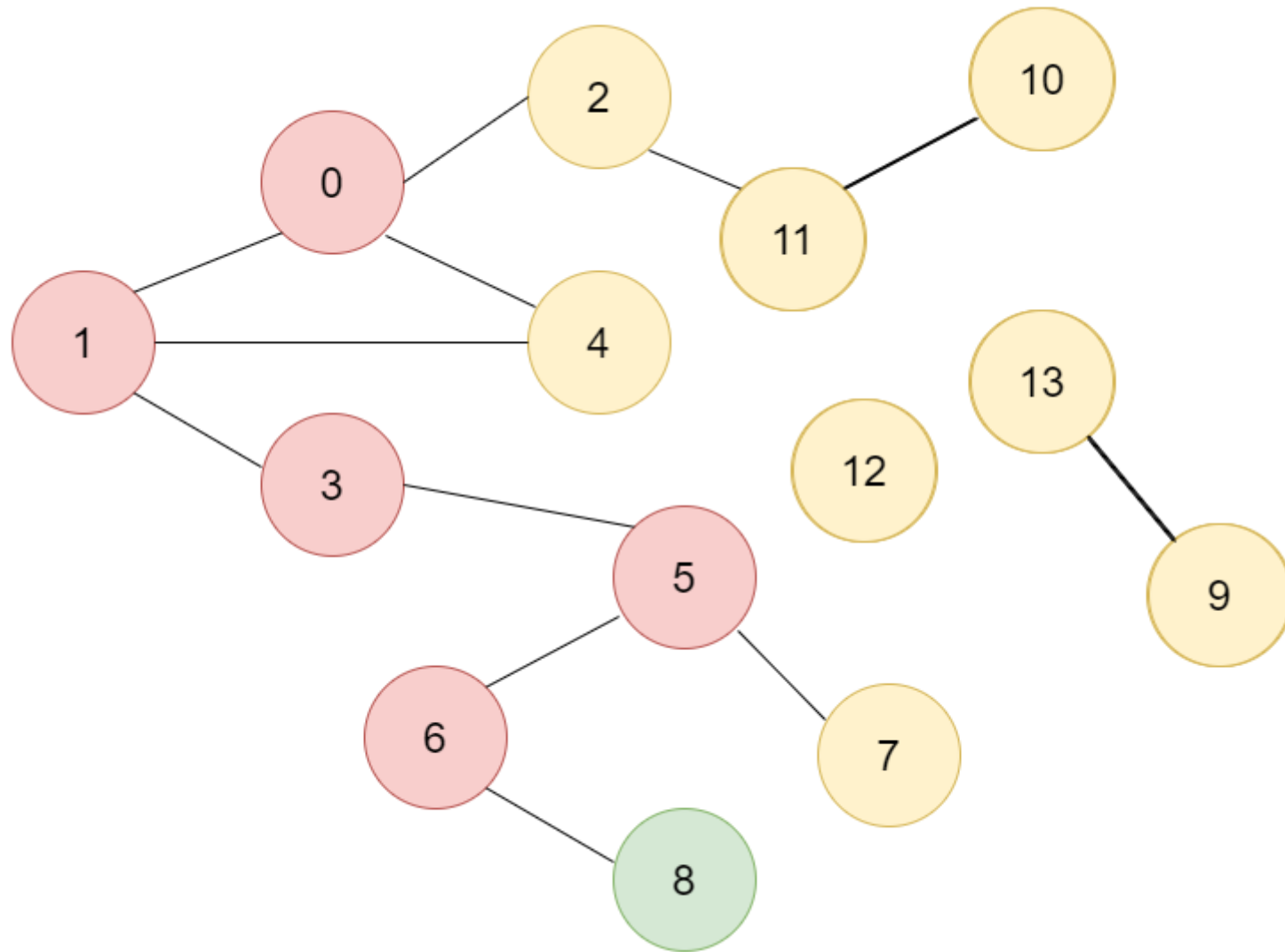


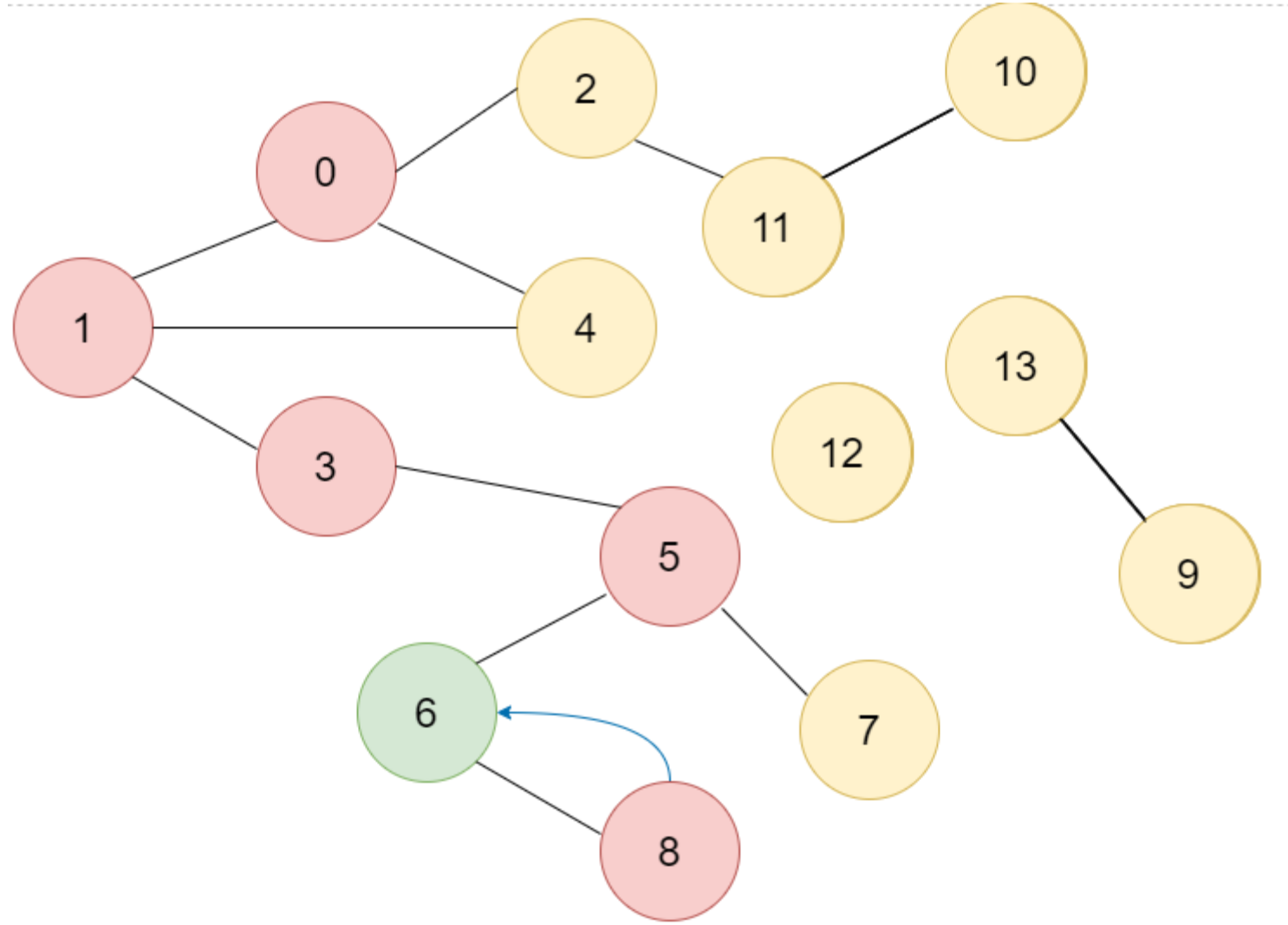


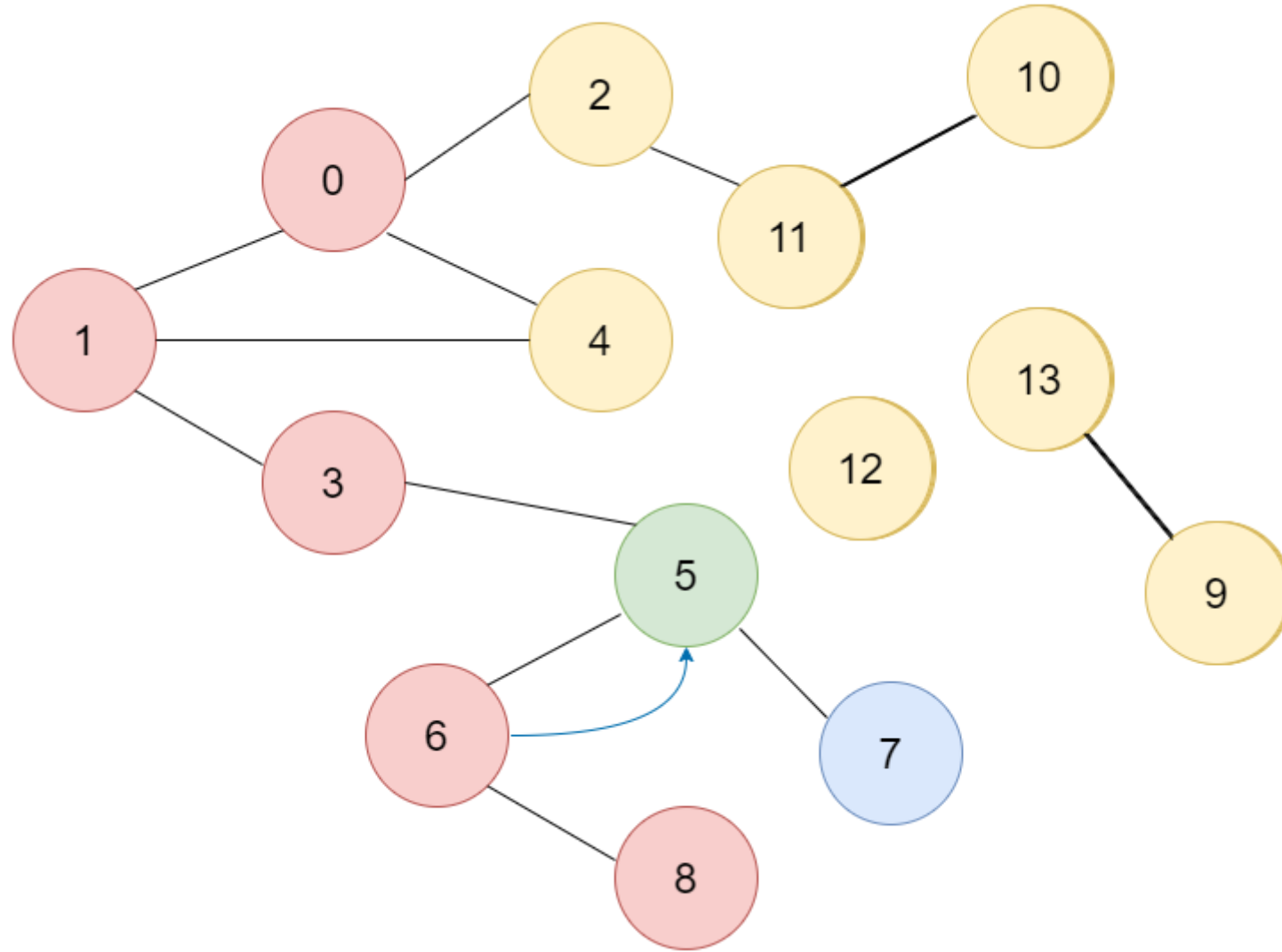


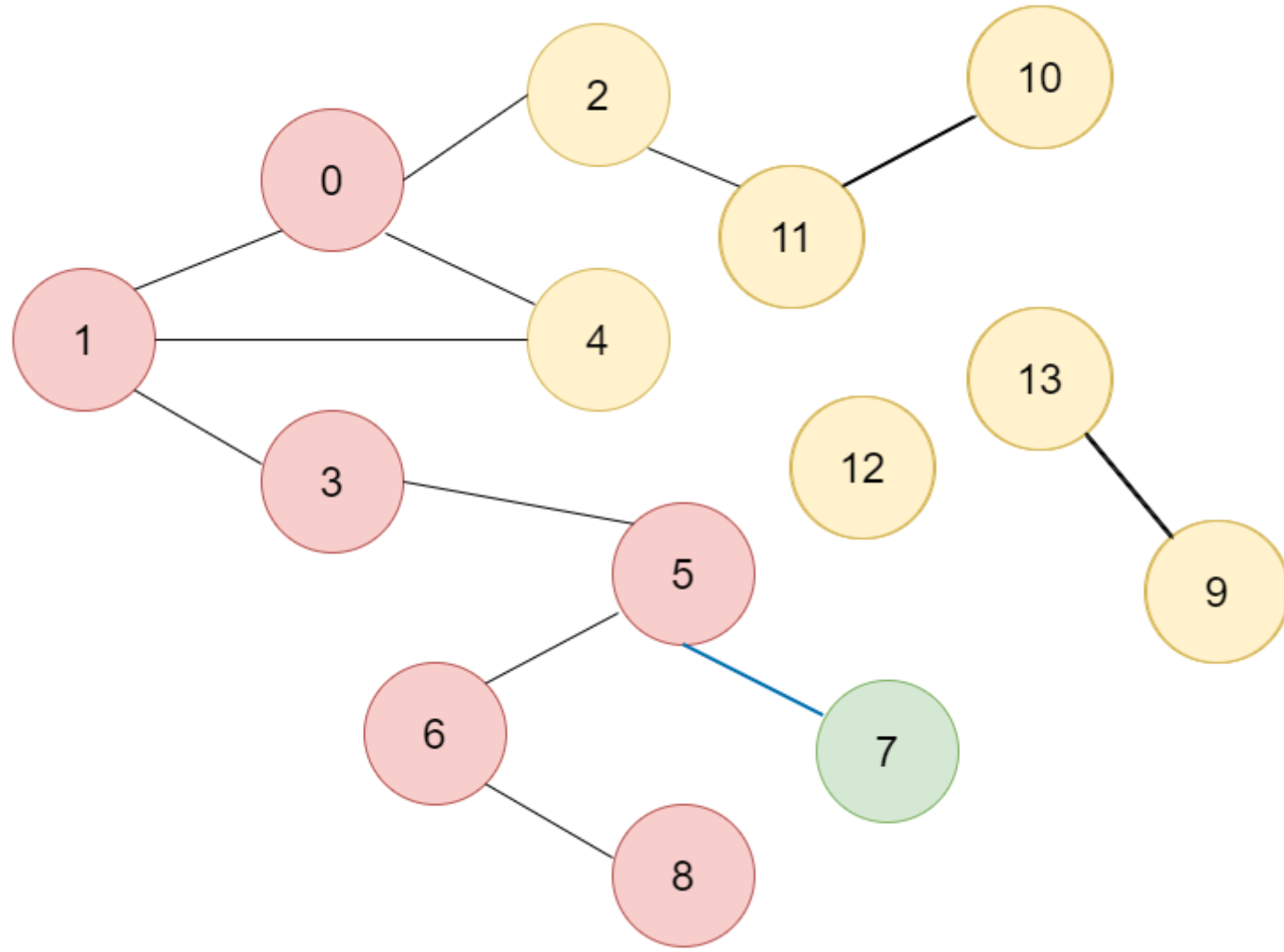


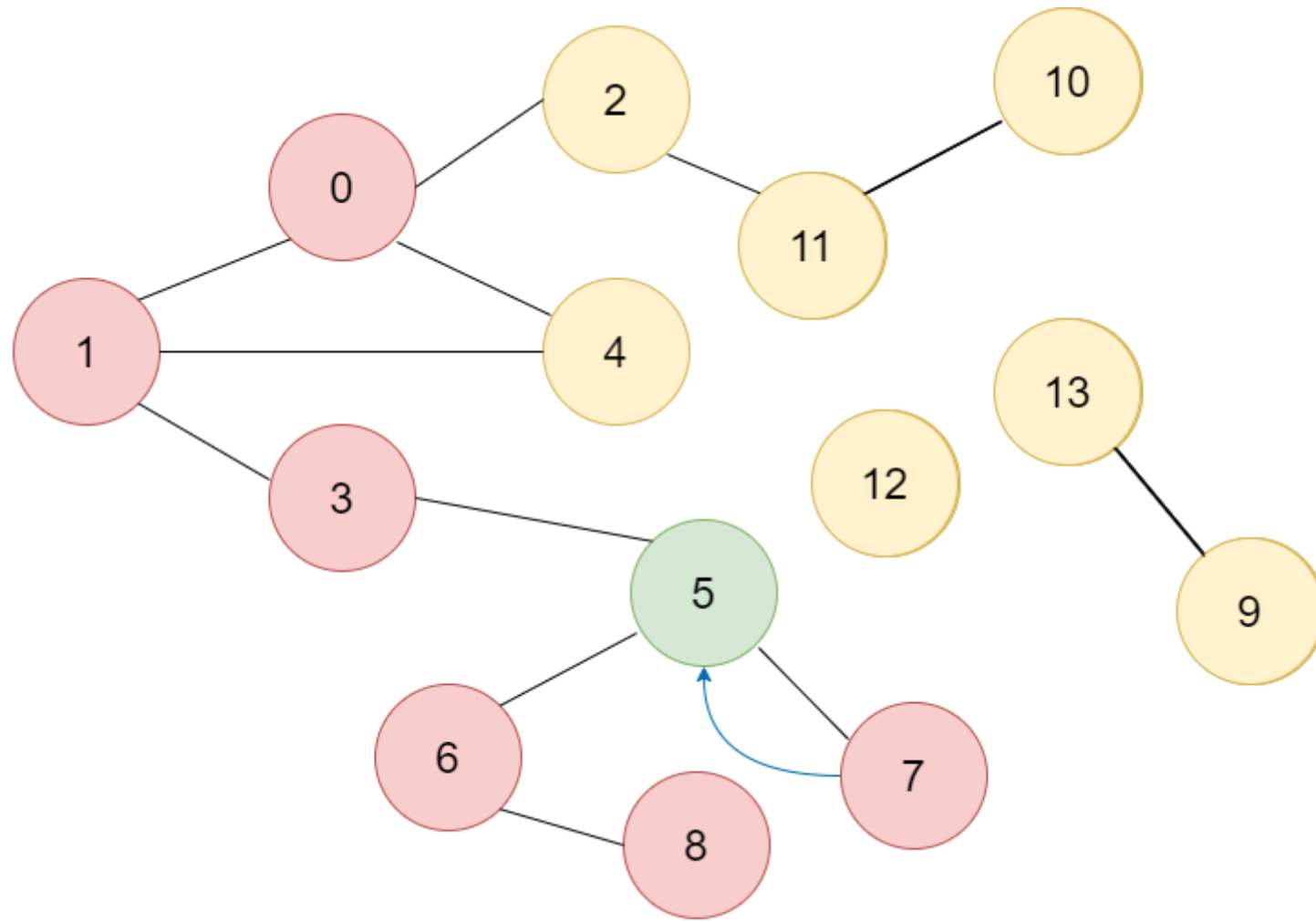


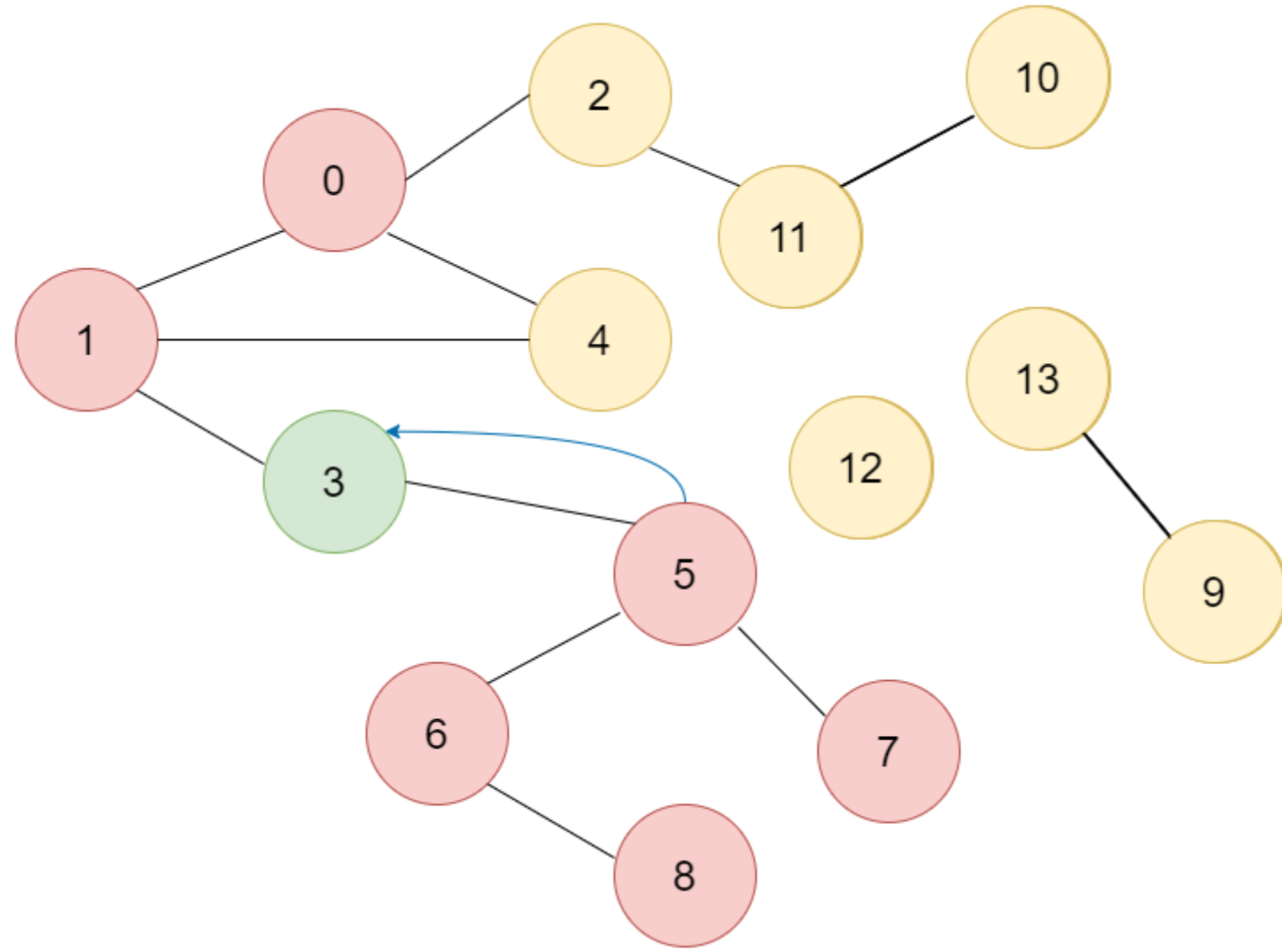




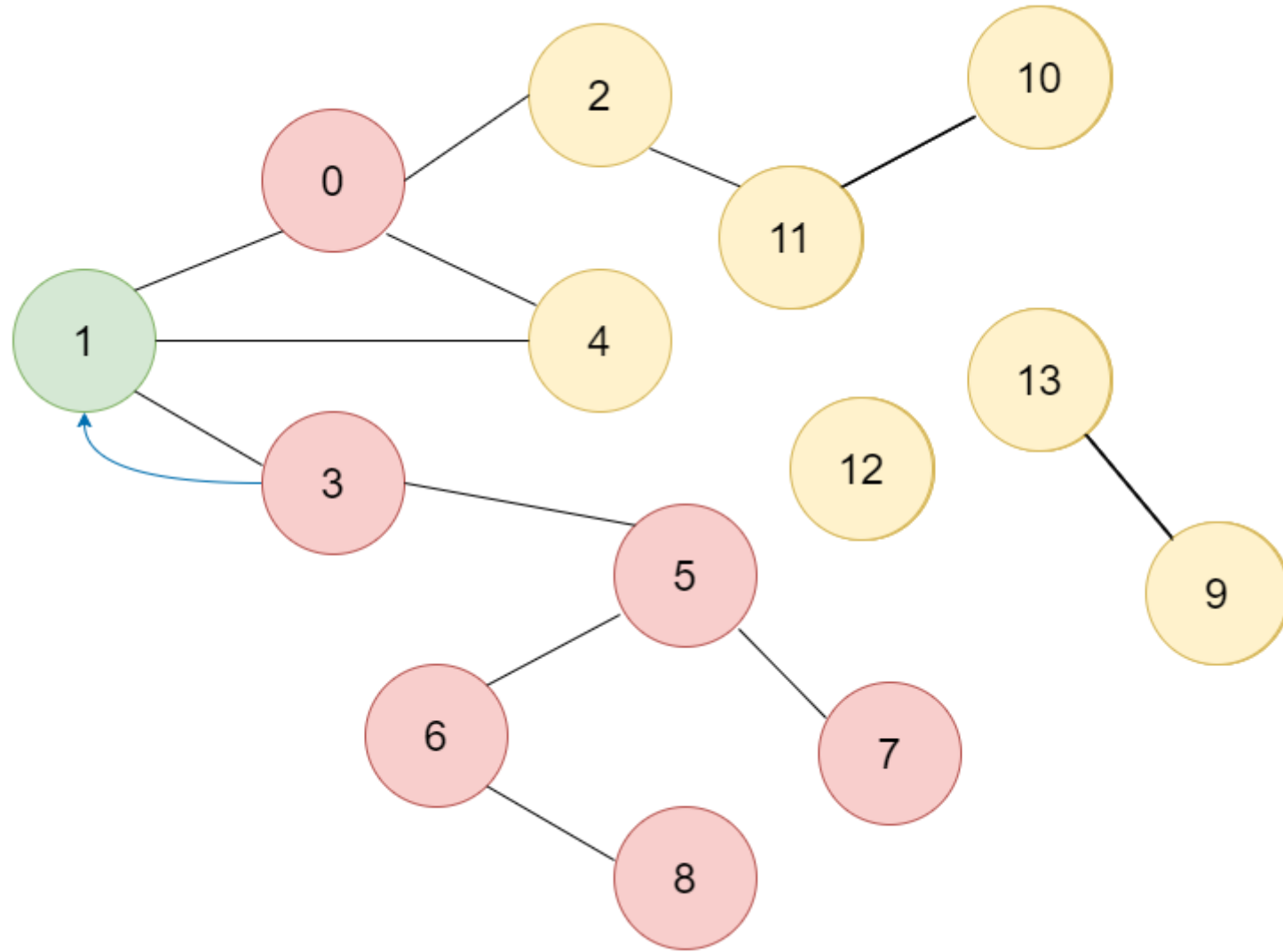


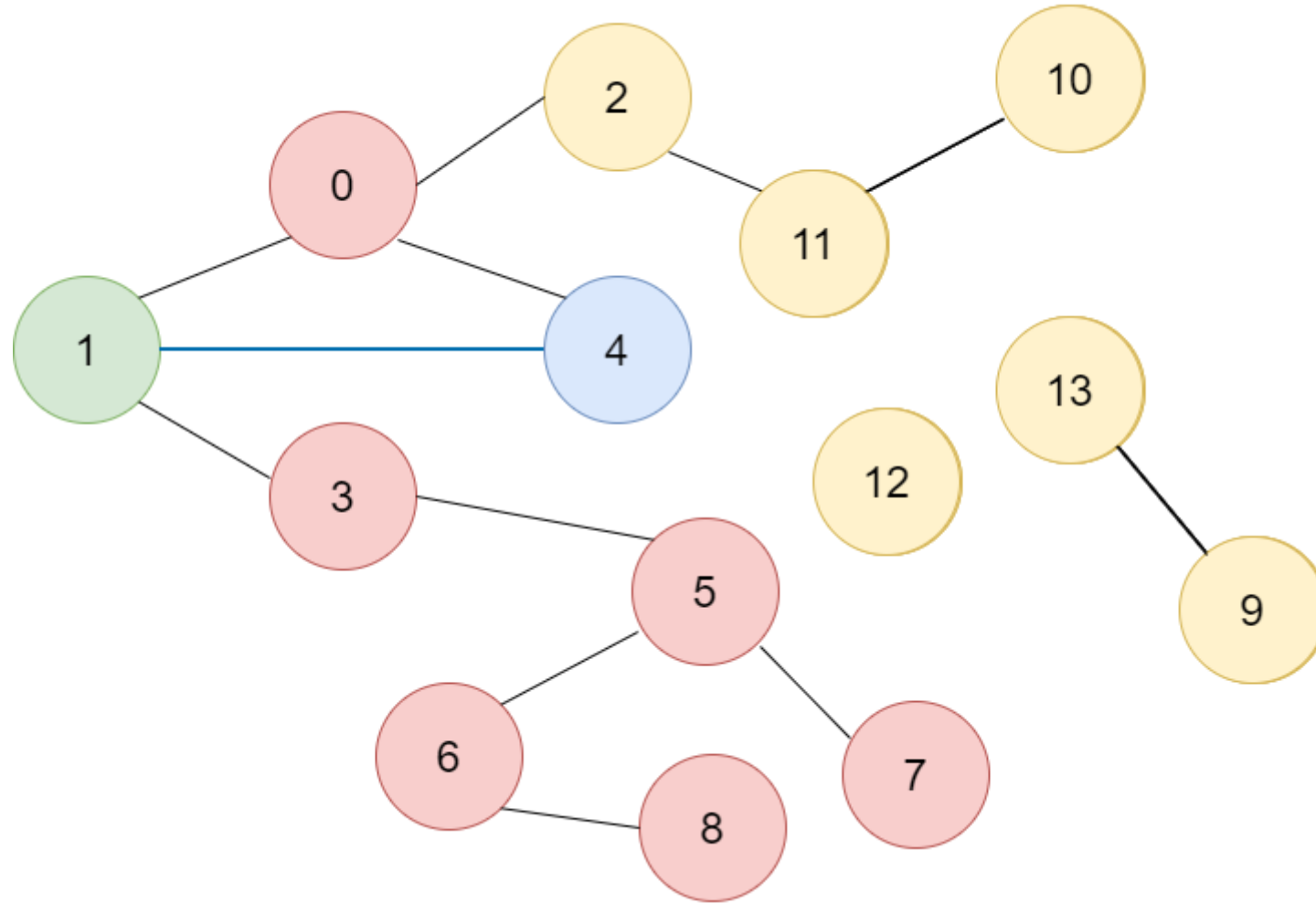


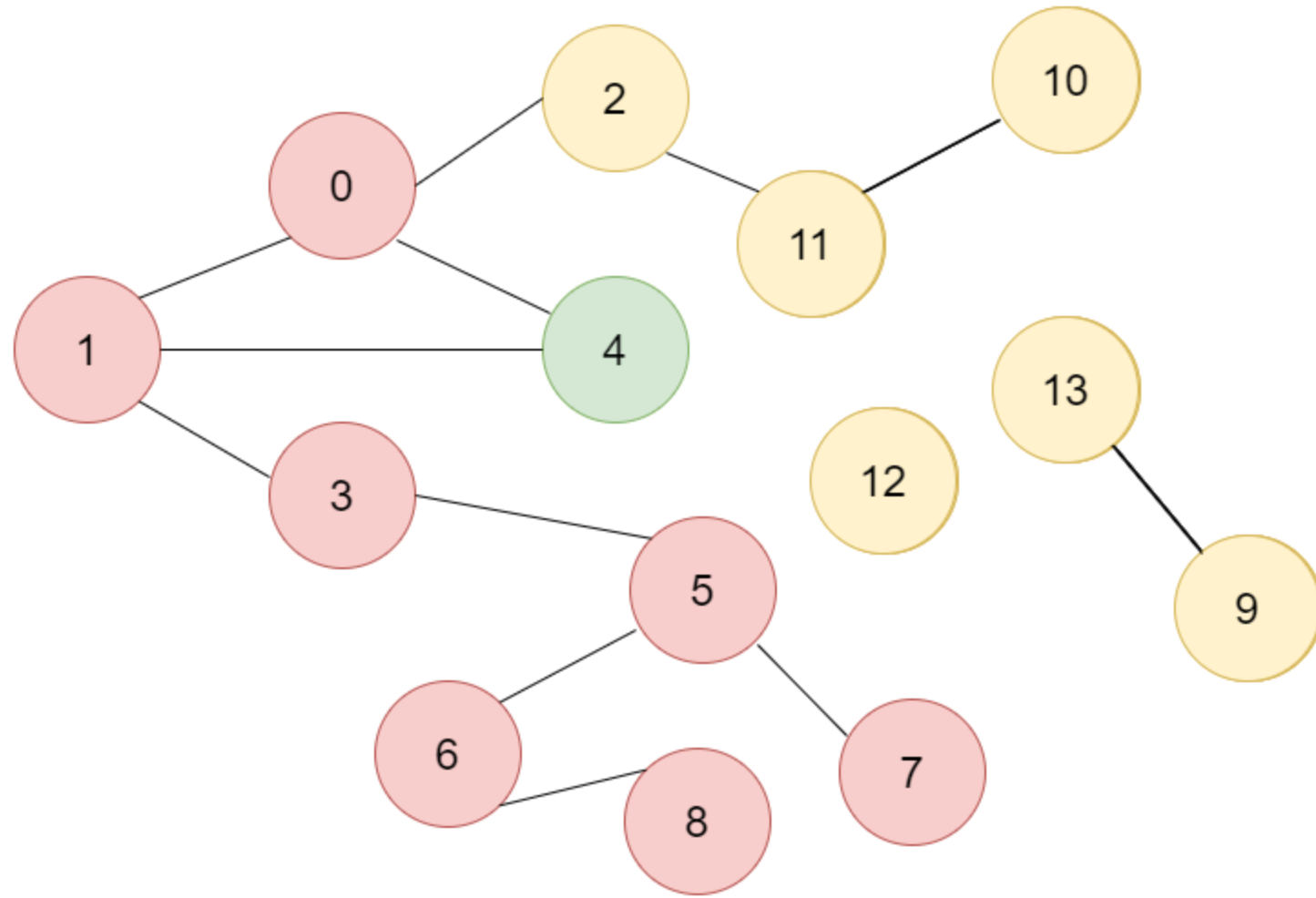


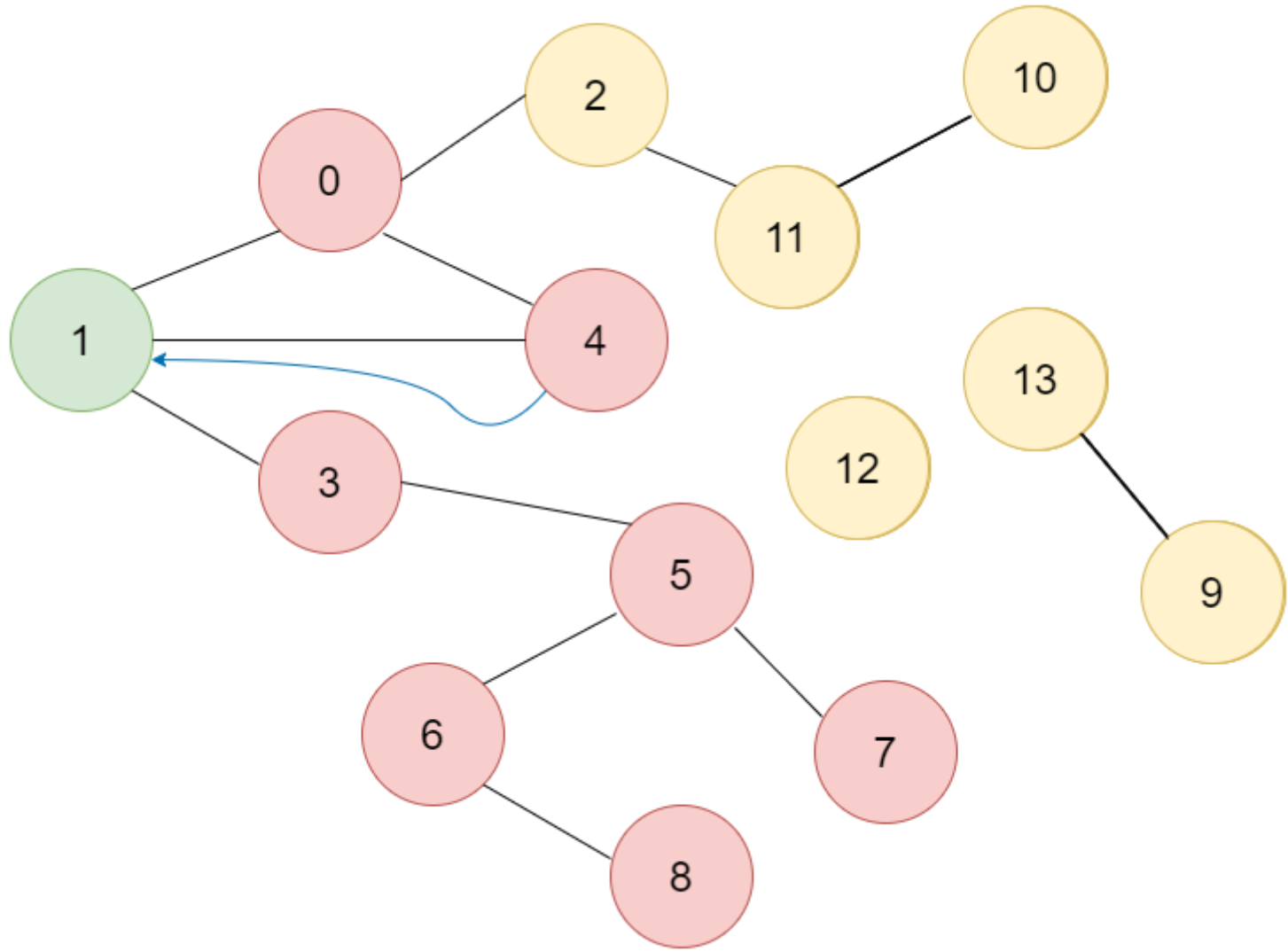


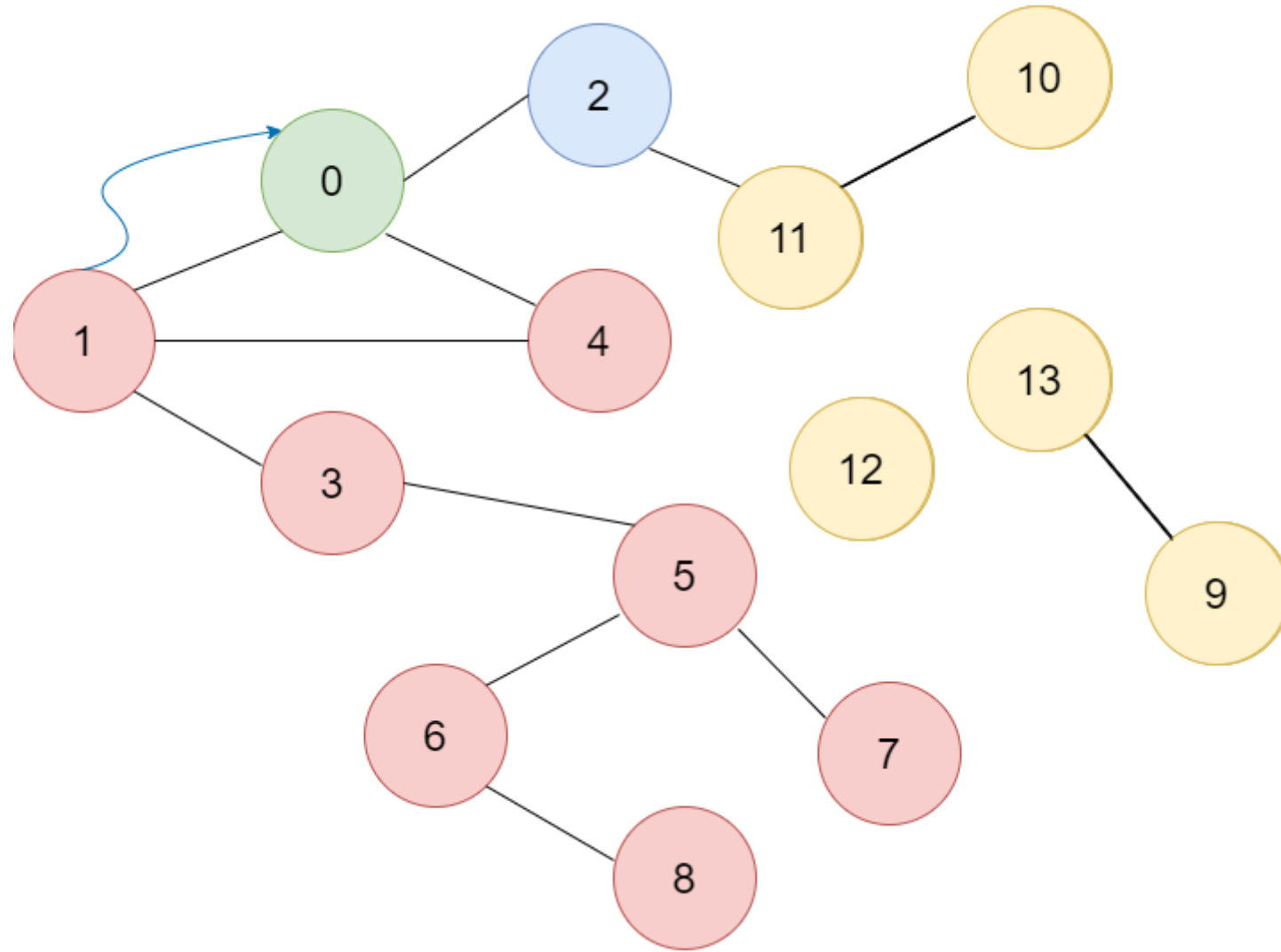


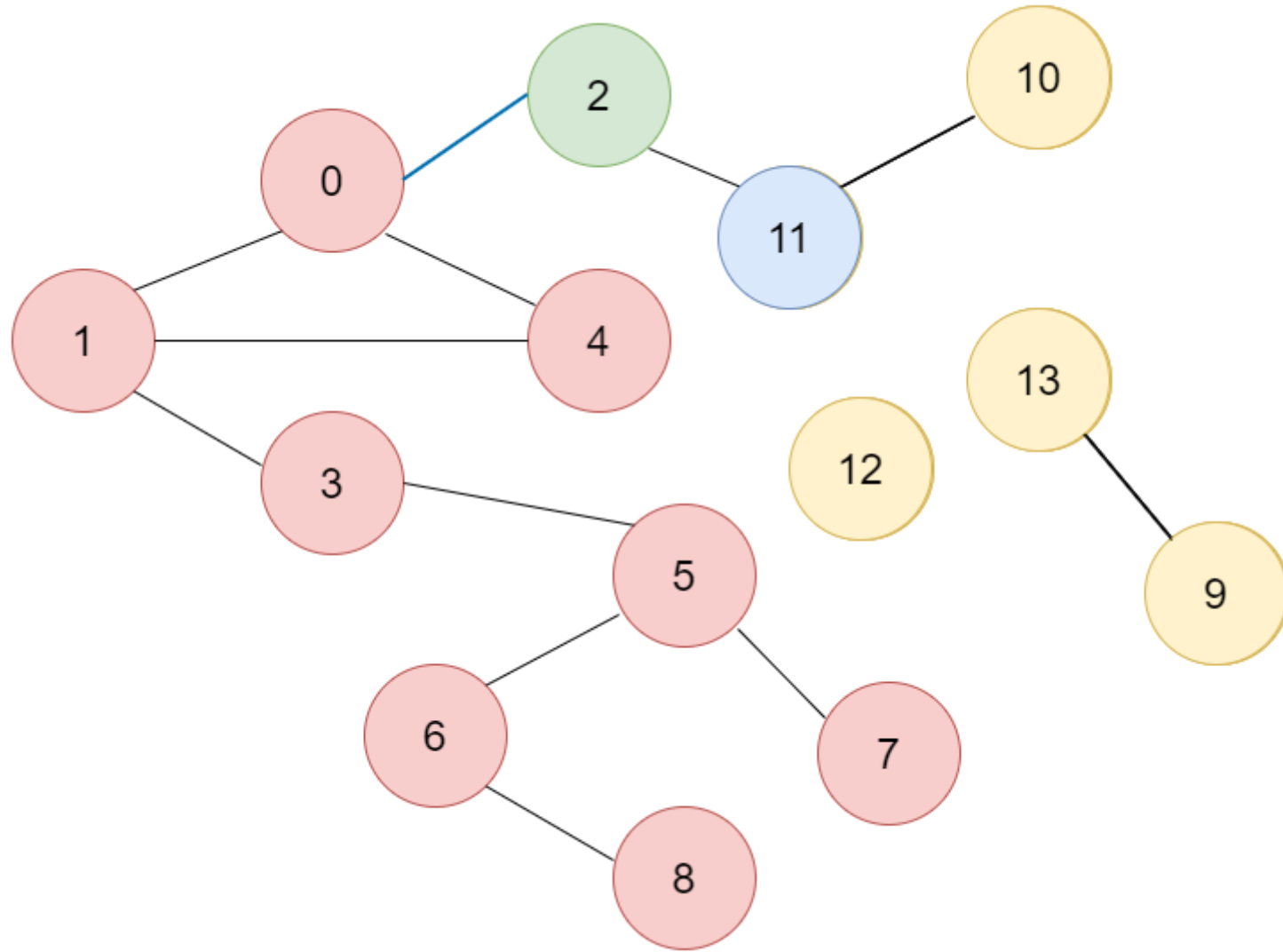


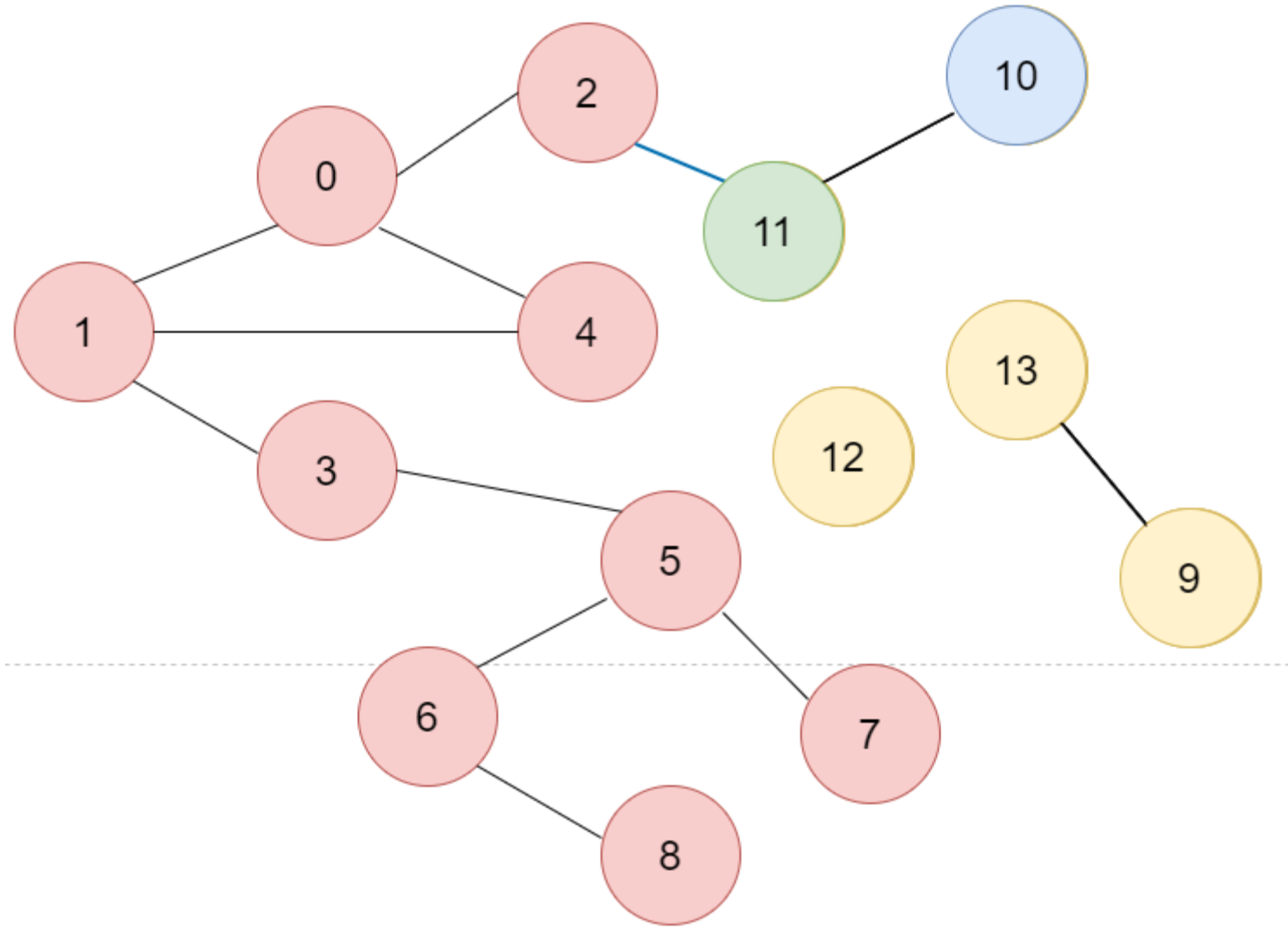


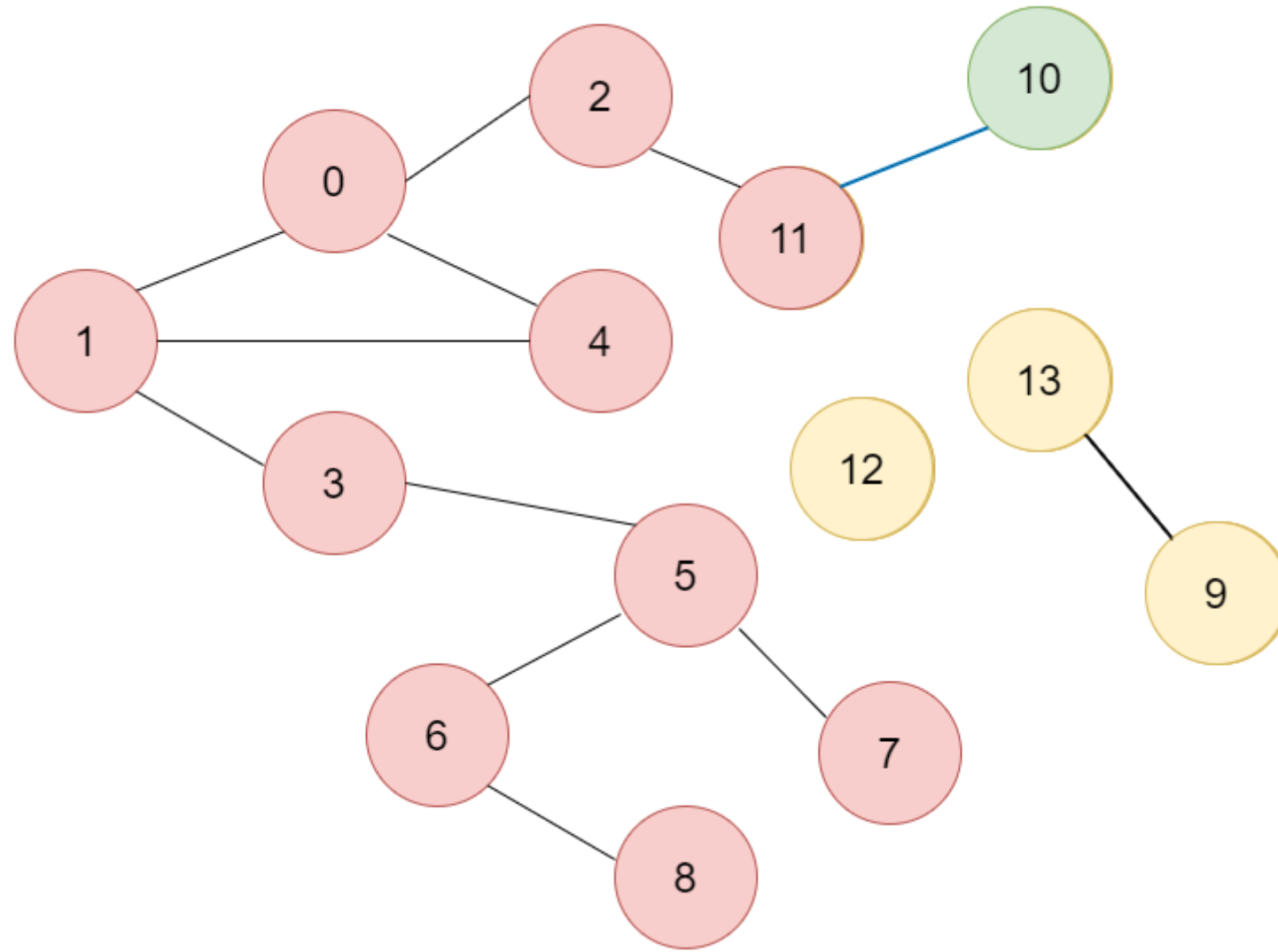




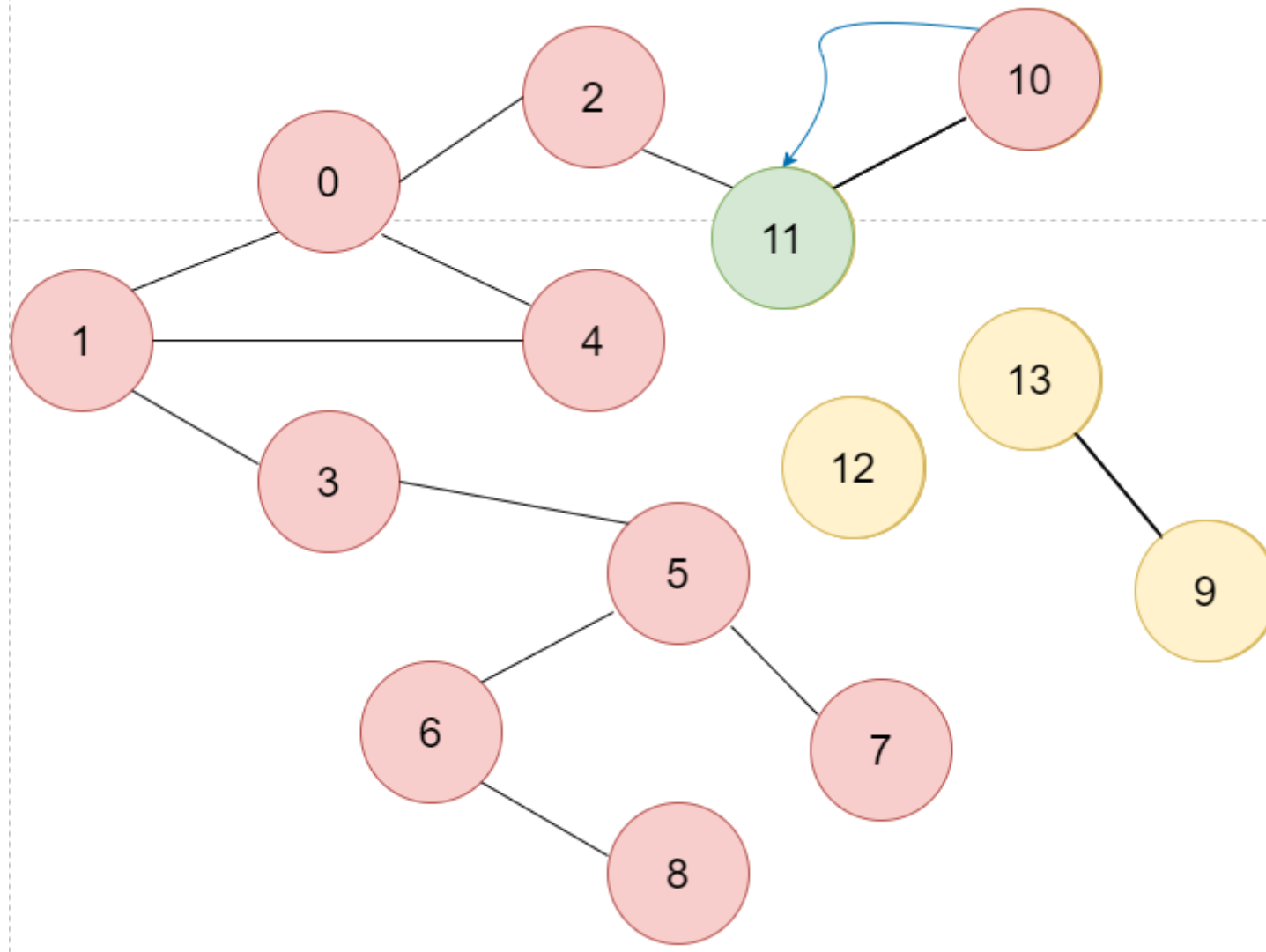


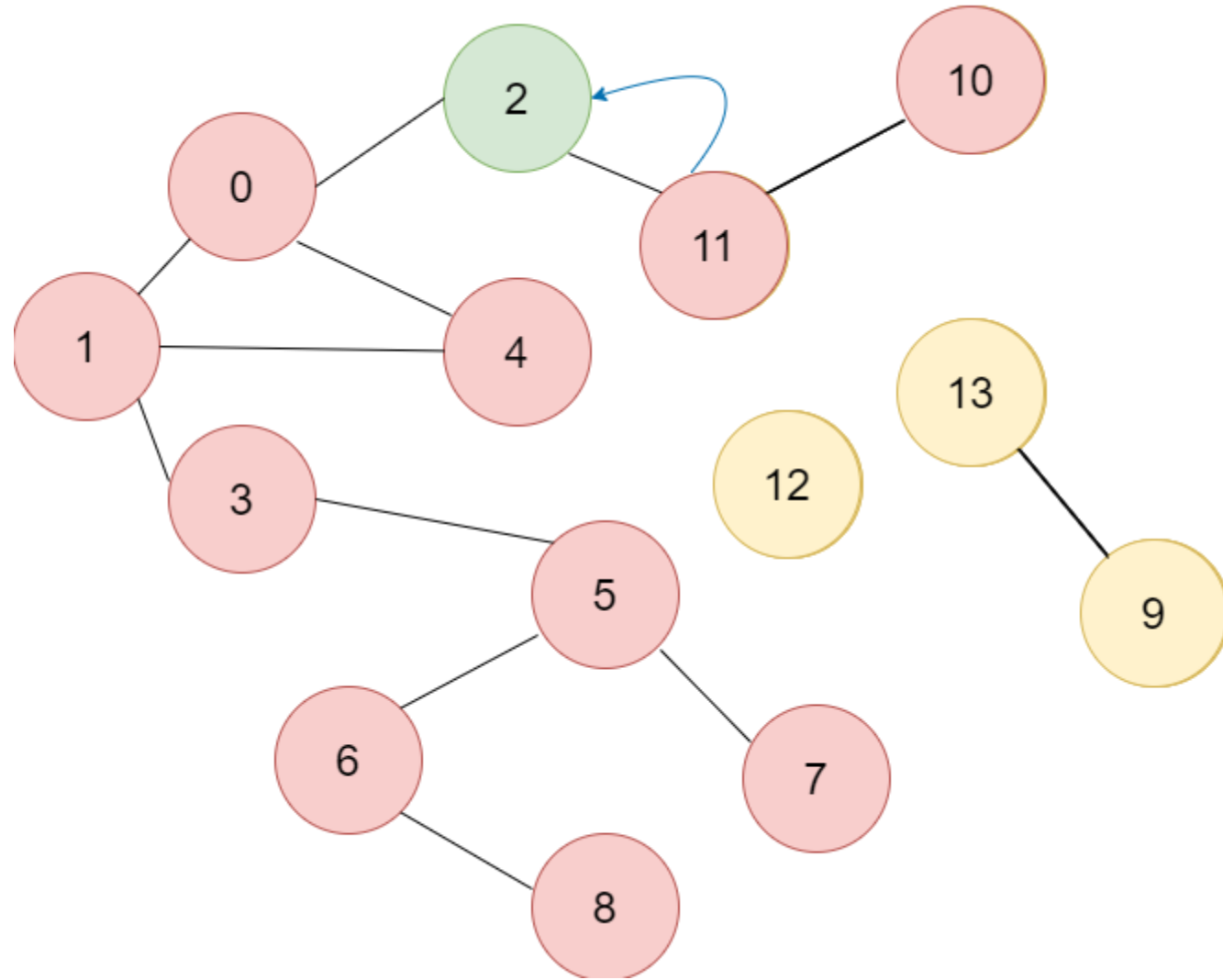


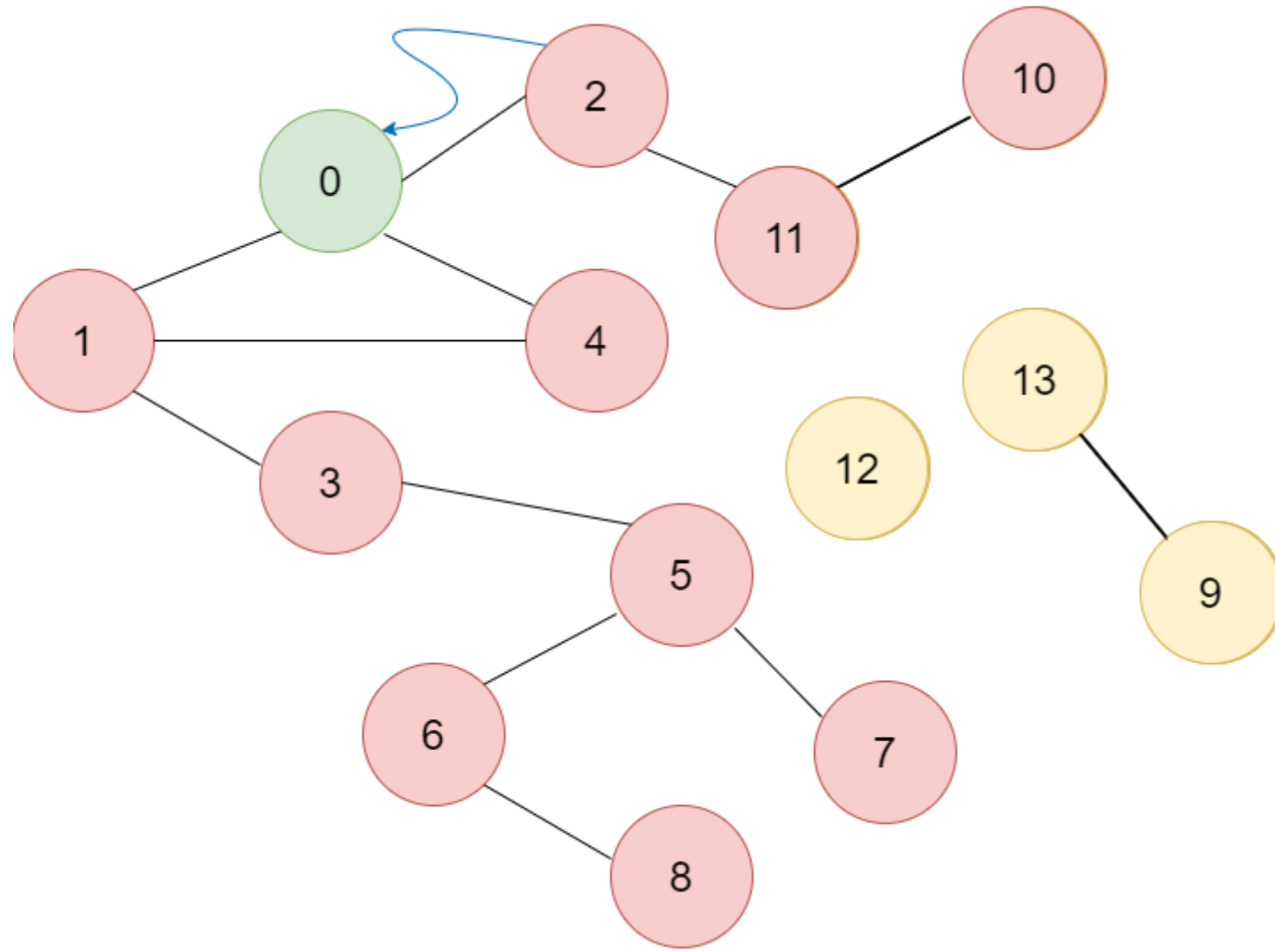


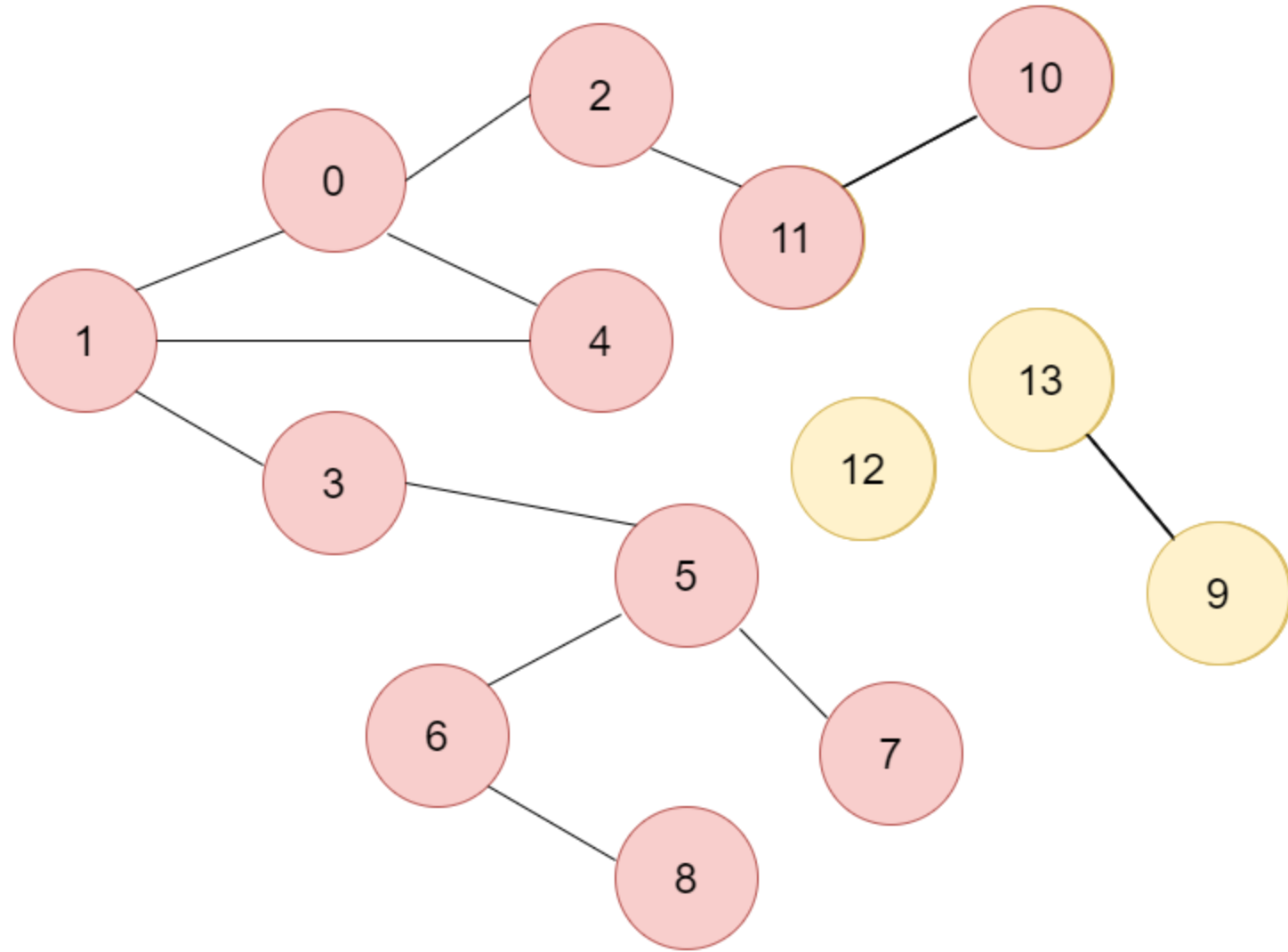












# PSEUDO CODE

**DFS (to visit a vertex  $v$ )**

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Mark  $v$  as visited.

Recursively visit all unmarked  
vertices  $w$  adjacent to  $v$ .

---

DFS( $u$ ):

```
Mark  $u$  as "Explored" and add  $u$  to  $R$ 
For each edge  $(u, v)$  incident to  $u$ 
    If  $v$  is not marked "Explored" then
        Recursively invoke DFS( $v$ )
    Endif
Endfor
```

---

Image source: KT Book

**DFS (Recursive Version)**

**Input:** graph  $G = (V, E)$  in adjacency-list representation, and a vertex  $s \in V$ .

**Postcondition:** a vertex is reachable from  $s$  if and only if it is marked as "explored."

---

```
// all vertices unexplored before outer call
mark  $s$  as explored
for each edge  $(s, v)$  in  $s$ 's adjacency list do
    if  $v$  is unexplored then
        DFS ( $G, v$ )
```

Image source: T. Roughgarden

Complexity:  $O(V + E)$

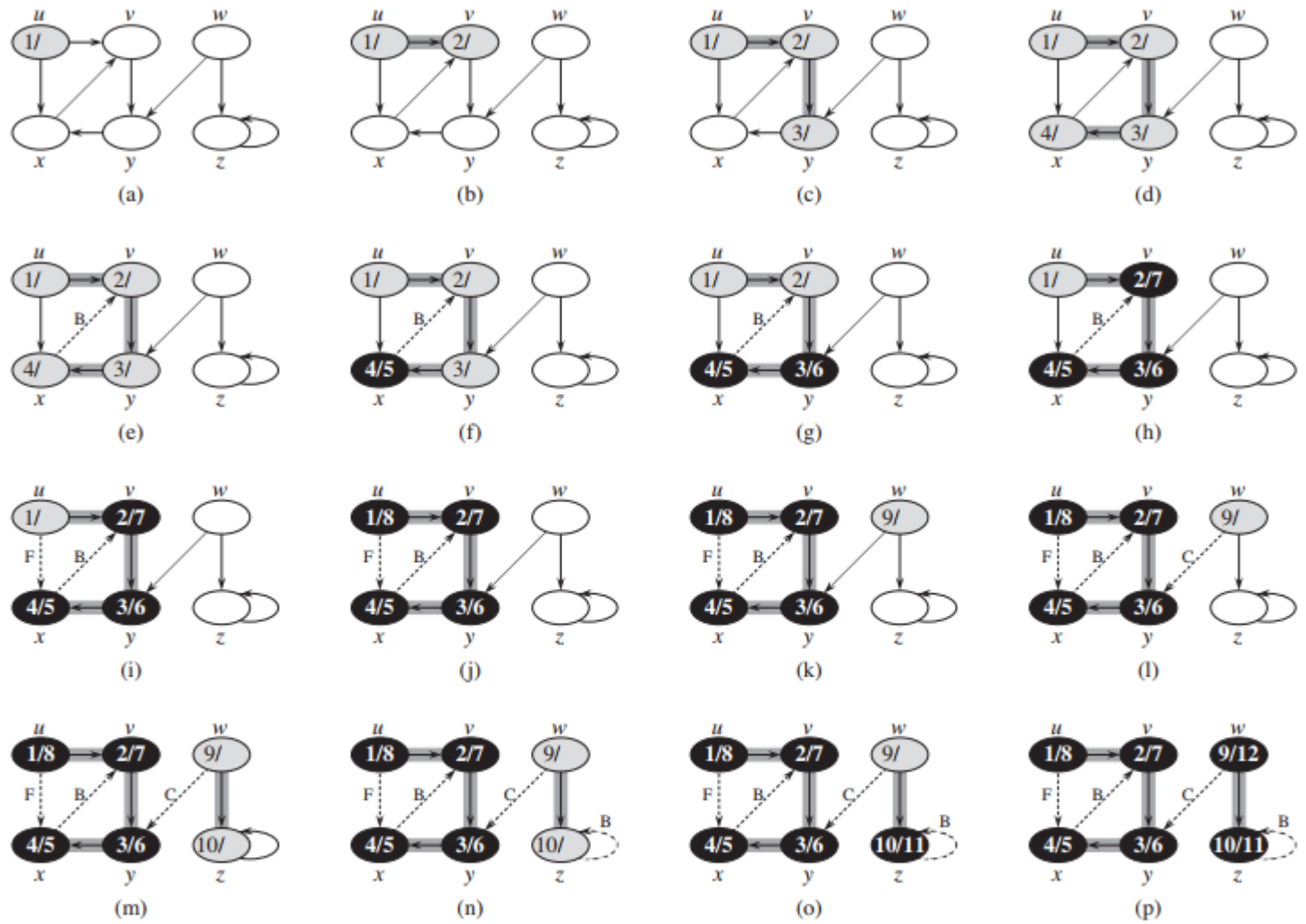


Image Source: CLRS

**Figure 22.4** The progress of the depth-first-search algorithm DFS on a directed graph. As edges are explored by the algorithm, they are shown as either shaded (if they are tree edges) or dashed (otherwise). Nontree edges are labeled B, C, or F according to whether they are back, cross, or forward edges. Timestamps within vertices indicate discovery time/finishing times.

DFS( $G$ )

```
1  for each vertex  $u \in G.V$ 
2       $u.color = \text{WHITE}$ 
3       $u.\pi = \text{NIL}$ 
4   $time = 0$ 
5  for each vertex  $u \in G.V$ 
6      if  $u.color == \text{WHITE}$ 
7          DFS-VISIT( $G, u$ )
```

DFS-VISIT( $G, u$ )

```
1   $time = time + 1$            // white vertex  $u$  has just been discovered
2   $u.d = time$ 
3   $u.color = \text{GRAY}$ 
4  for each  $v \in G.Adj[u]$      // explore edge  $(u, v)$ 
5      if  $v.color == \text{WHITE}$ 
6           $v.\pi = u$ 
7          DFS-VISIT( $G, v$ )
8   $u.color = \text{BLACK}$        // blacken  $u$ ; it is finished
9   $time = time + 1$ 
10  $u.f = time$ 
```

NEXT TOPIC?

Depth First Search (DFS) with Stack